

Security Challenges in Virtualized Environments

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Invisible Things Lab

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San Francisco, April 8th 2008

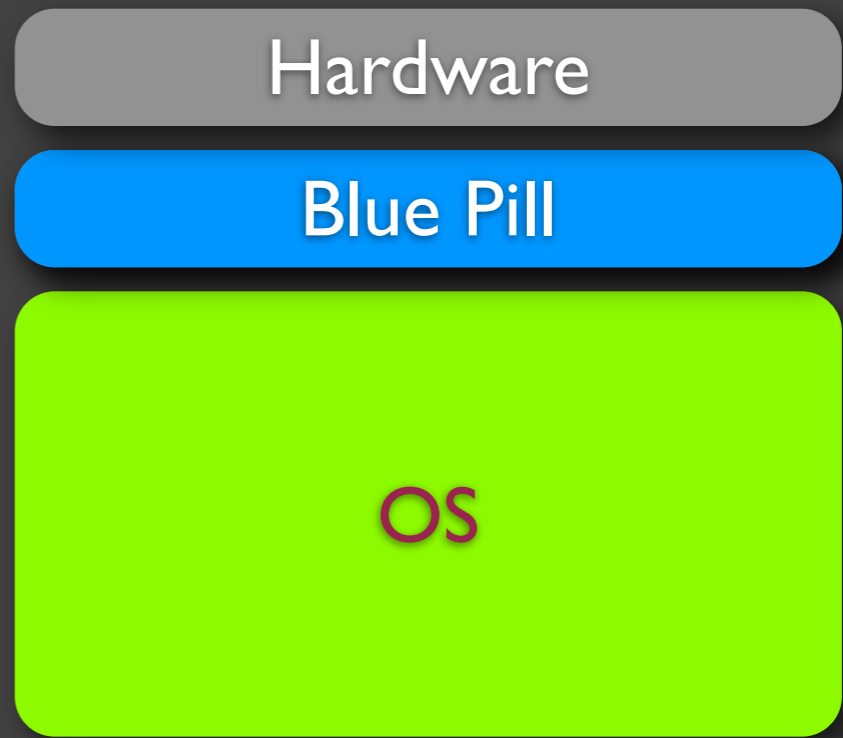
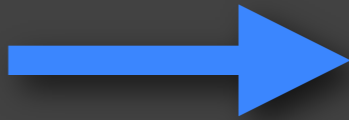
1 Virtualization-based **MALWARE**

2 Using Virtual Machines for **ISOLATION**

3 **NESTED** virtualization



Virtualization-based MALWARE



AMD-V
Intel VTx

Blue Pill Characteristics

NO HOOKS!



Cannot be detected using any integrity scanner

On the fly installation



No boot/BIOS/etc modifications necessary

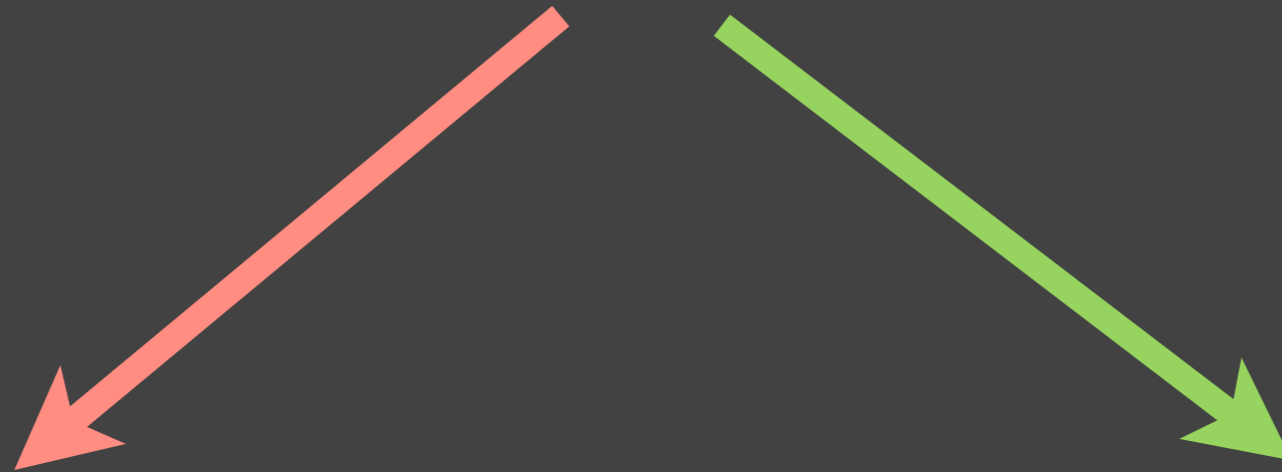
No I/O virtualization



Negligible performance impact (your brand new 3D card will still work!)

Blue Pill detection

Blue Pill detection



Detecting a VMM

Detecting
virtualization based
malware

VMM detection

Direct timing analysis

Guest time virtualization

HPET timers

Blue Chicken

CPU specific behavior

TLB profiling

VMM detection?

- Everything is going to be virtualized!
- Thus the information that “there is a hypervisor in the system”...
- ...would be pretty much useless...

Detecting virtualized
malware?

No Hooks!

Search for code

Detect activity
(e.g. network packets)

Heuristics

By Pattern

- Stealth by Design concept
- Covert channels

Nested Page Tables
(hardware SPT)

“Massive” malware

Simple
Obfuscation

Won't work

0day malware

But why we can't use obfuscation for "classic" malware?

Because it leaves **hooks** anyways!

And **we can always find those hooks**, no matter how obfuscated the classic malware is!

The whole big deal about Blue Pill is:

NO HOOKS in the *system!*

Blue Pill prevention

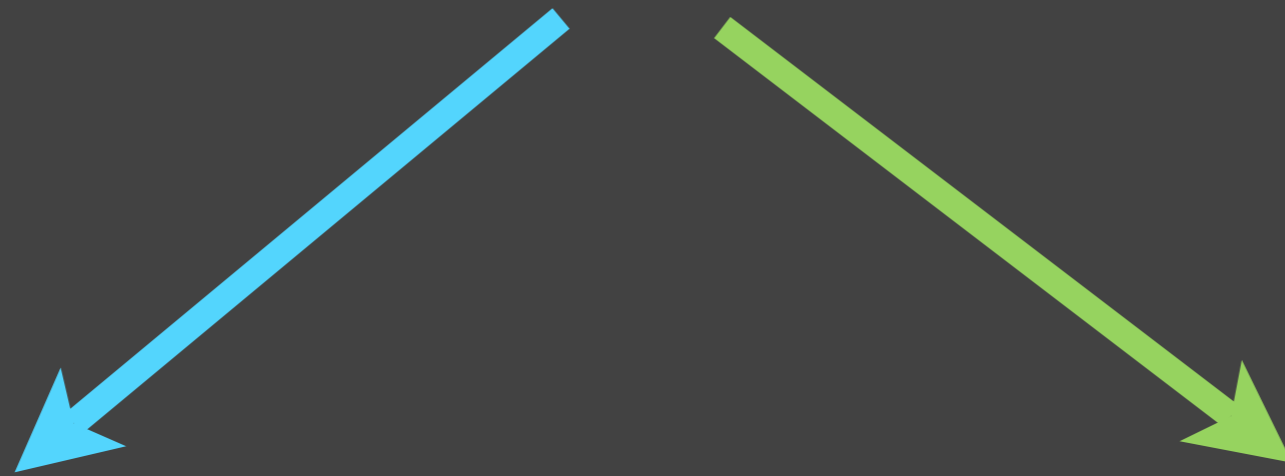


Disable virtualization?

How about also disabling your network card so you never got infected from the Internet?

Install a **trusted**
hypervisor first?

Installing trusted hypervisor



Static Root of Trust
Measurement

BIOS > MBR > VMM
e.g. MS Bitlocker

Dynamic Root of Trust
Measurement

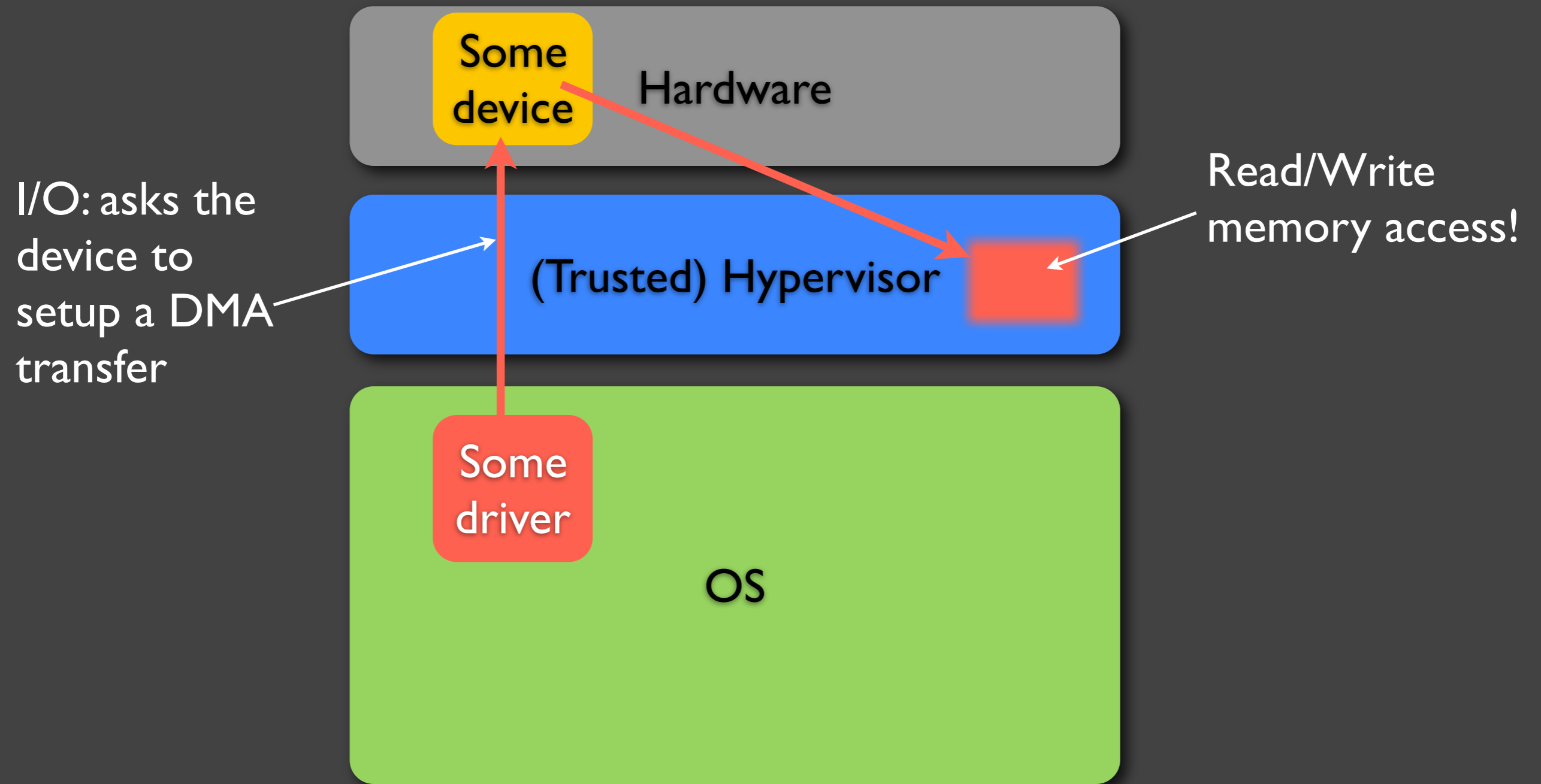
SENTER (Intel TXT)
SKINIT (AMD SVM)

Trusted vs. Secure?

- SRTM and DRTM only assures that what we **load** is **trusted**...
- ...at the moment of **loading**!
- 3 sec later... it could be exploited and get compromised!

Trusted \neq Secure (e.g. flawless)

E.g. #1: The famous DMA problem



IOMMU

- Solution to the problem of “DMA attacks”
- Intel calls it: VT-d
- Not much PC hardware supports it yet
 - Expected to change soon
- No THIN HYPERVISORS without IOMMU!

Other problems with VMMs?
Stay tuned...

All in all: it's not trivial to have a trusted & secure
hypervisor installed...
... but for sure this is the proper way to go...

Virtualization-based **MALWARE**

2 Using Virtual Machines for **ISOLATION**

3 **NESTED** virtualization



Using Virtual Machines for ISOLATION

Originally **ISOLATION** was supposed to be provided by
Operating Systems...

- Separate processes/address spaces,
- User accounts & ACLs...

But in practice current OSes simply
fail at providing isolation!

Why OSes fail?

- Kernel bugs!
- Kernel bugs!!
- Kernel bugs!!!
- Bad design, e.g.:
 - XP and “all runs as admin” assumption
 - Vista’s UAC assumes admin rights should be granted to every installer program!

VMMs for the rescue!

trusted & secure hypervisor

Vista
(work projects)

Linux + Firefox
("random"
surfing)

Linux + Firefox
(online banking)

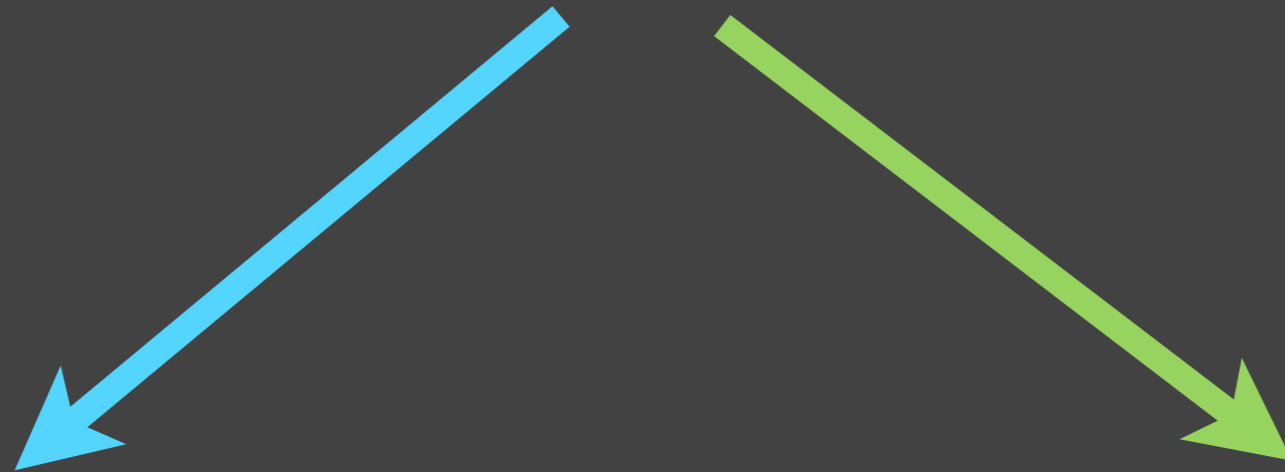
MacOSX
("home", e.g.
pics, music, etc)

Challenges

- Performance
- Why is VMM/hypervisor going to be more secure than OS's kernel?

VMM bugs?

VMM Bugs



Bugs in hypervisors

Bugs in additional
infrastructure

E.g. #1: CVE-2007-4496

- VMWare ESX 3.0.1
 - <http://www.vmware.com/support/vi3/doc/esx-8258730-patch.html>
- Found by Rafal Wojtczuk (McAfee)
- September 2007
- Guest OS can cause memory corruption on the host and *potentially* allow for arbitrary code execution on the host

E.g. #2: CVE-2007-0948

- Microsoft Virtual Server 2005 R2
 - <http://www.microsoft.com/technet/security/bulletin/ms07-049.msp>
- Found by Rafal Wojtczuk (McAfee)
- August 2007
- Heap-based buffer overflow allows guest OS to execute arbitrary code on the host OS

E.g. #3: CVE-2007-4993

- Xen 3.0.3
 - http://bugzilla.xensource.com/bugzilla/show_bug.cgi?id=1068
- Found by Joris van Rantwijk
- September 2007
- By crafting a grub.conf file, the root user in a guest domain can trigger execution of arbitrary Python code in domain 0.

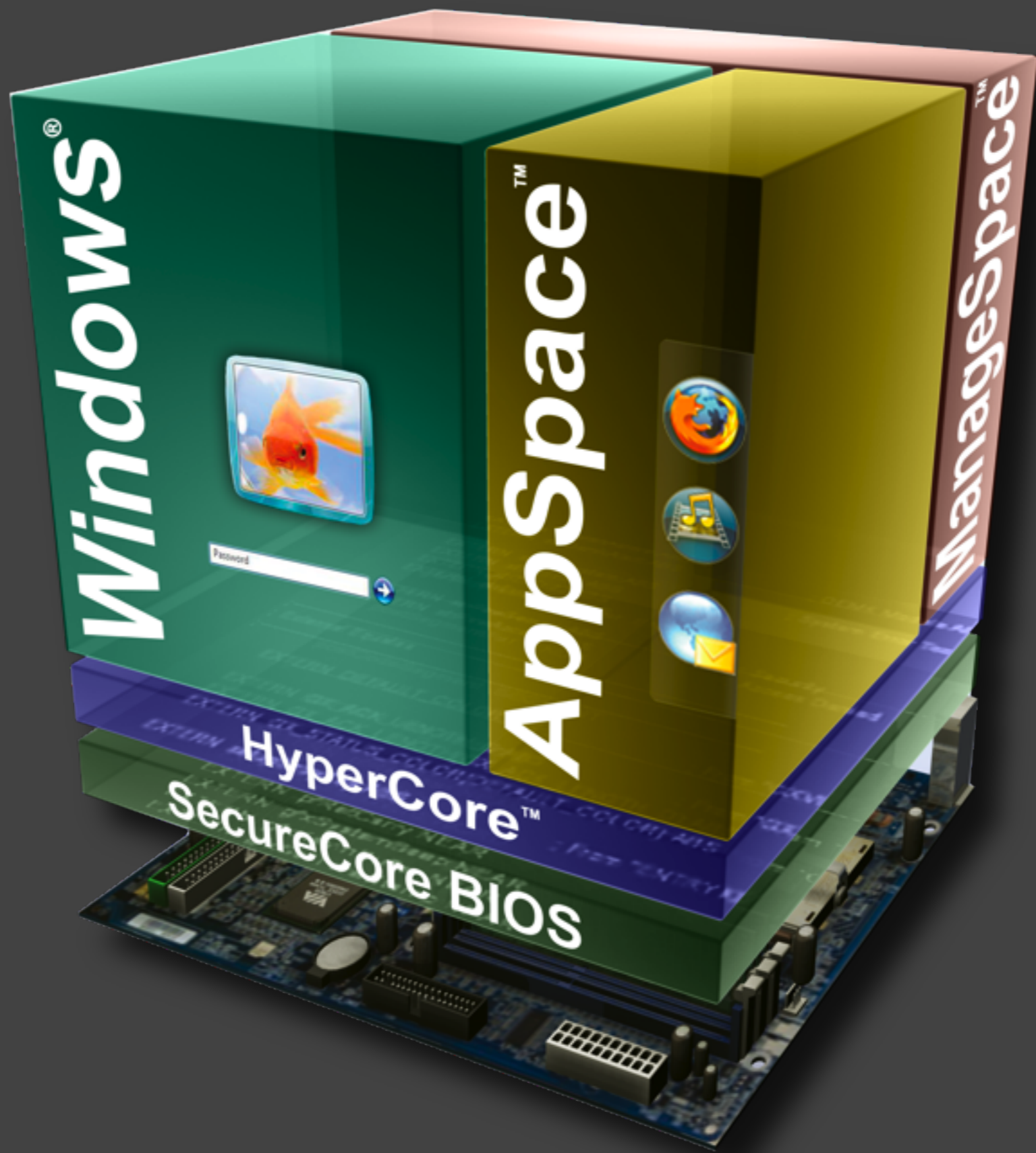
E.g. #4: Various Bugs

- Paper by Tavis Ormandy (Google)
 - <http://taviso.decsystem.org/virtsec.pdf>
- April 2007
- Disclosed bugs in VMWare, XEN, Bochs, Virtual PC, Prallels
- A simple fuzzers for:
 - Instruction parsing by VMMs
 - I/O device emulation by VMMs

As you see current VMMs are far from being flawless...

To make VMMs more secure we need to keep them
ultra-thin and small!

Phoenix HyperSpace



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e.g., "10 market st, san francisco" or "hotels near lax"

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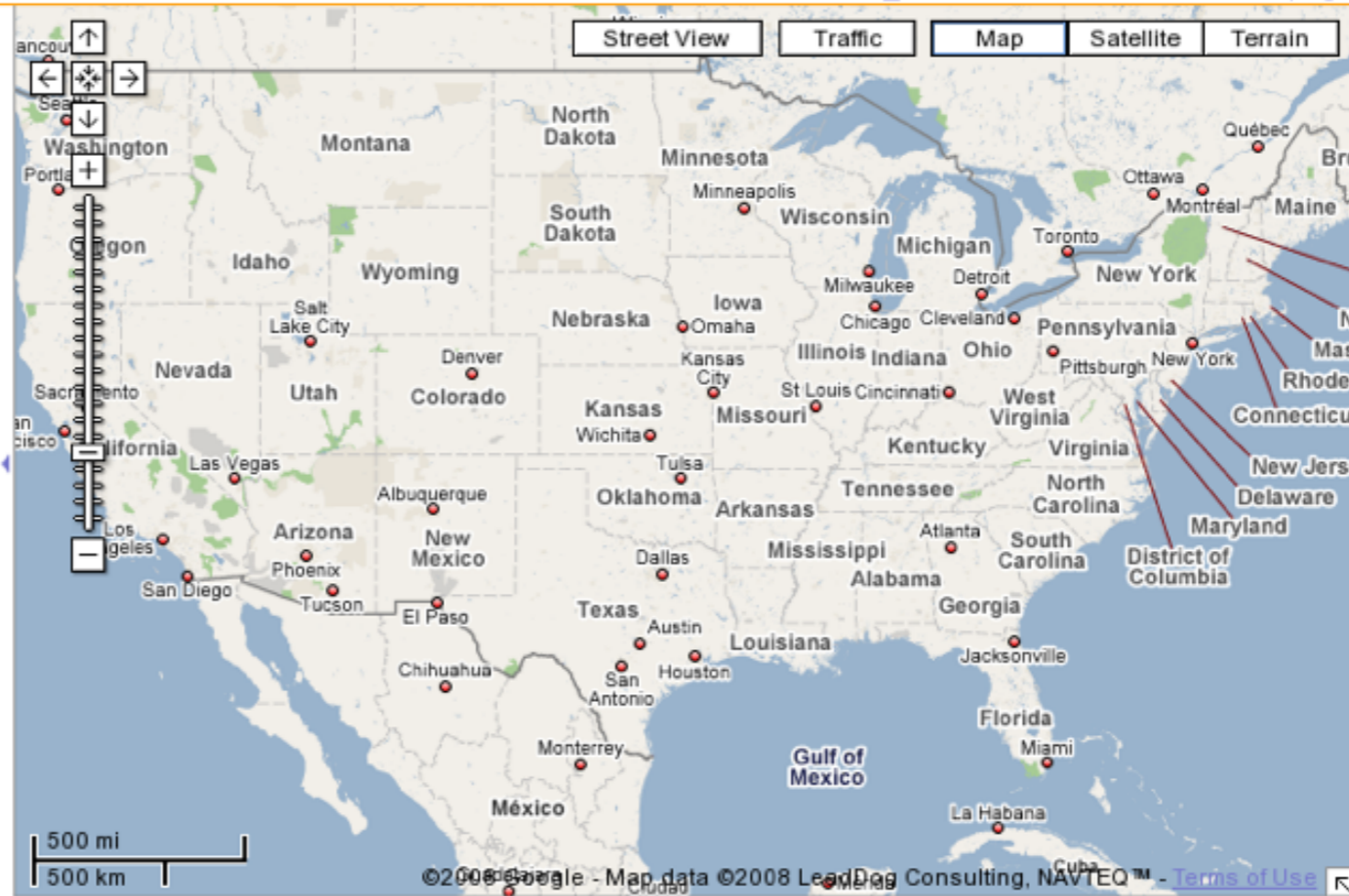
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Done

Two Kennedys endor...
5 U.S. troops killed...
Burglar steals bishop...
Militants release Pak...

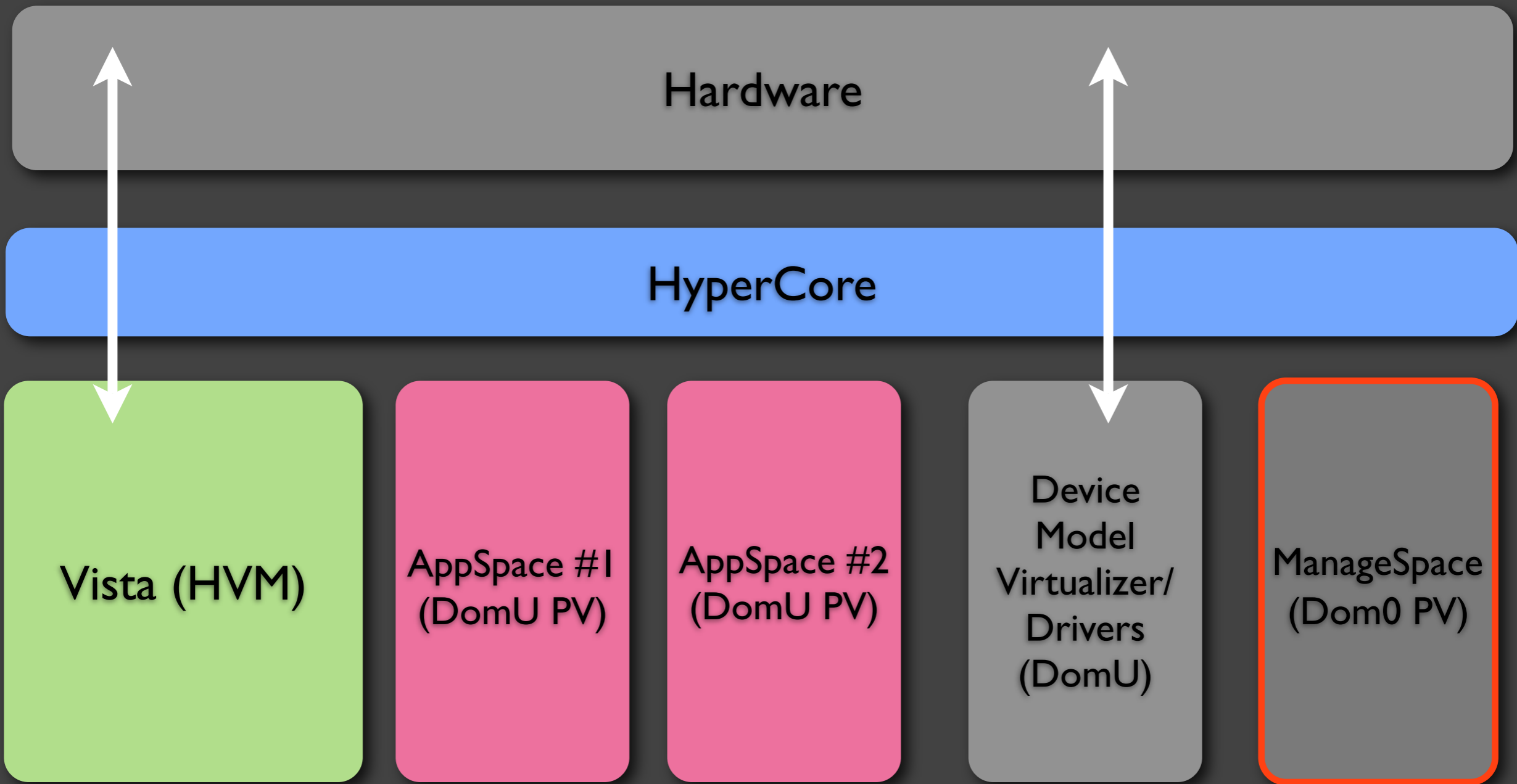


Recycle Bin
cmd - Shortcut



2:39 PM

HyperCore:
the type I hypervisor used for HyperSpace



Hardware

HyperCore

Vista (HVM)

AppSpace #1
(DomU PV)

AppSpace #2
(DomU PV)

Device
Model
Virtualizer/
Drivers
(DomU)

ManageSpace
(Dom0 PV)

The HyperCore

- Targets **desktop/laptop systems**
- Guest OS execute at near-native performance (**including fancy graphics**)
- Support for full **ACPI** (Power Management)
- Integrity: loaded via SecureCore BIOS (**Static Root of Trust Measurement**)
- Very thin - easy to audit!

Speeding things up

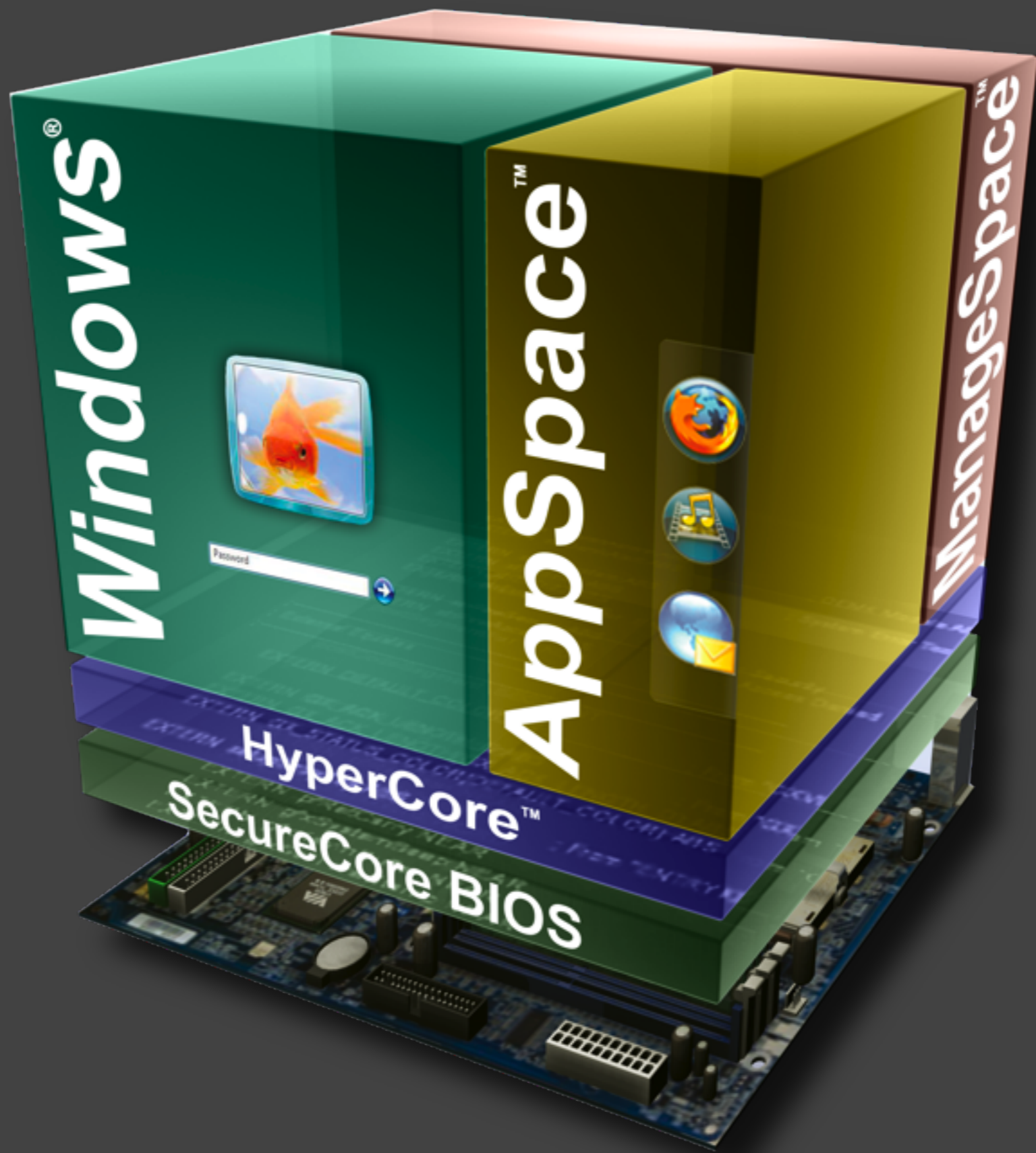
- Pass through for most devices
- SPT: 1-1 mapping for most pages for the Primary OS

Power Management

- ACPI tables exposed to the Primary OS, so that the overall power performance is optimized
- Efficient intercepts for power management control

Integrity

- Static RTM via Phoenix's SecureCore BIOS
- Dynamic RTM via Intel's TXT/AMD's SKINIT
- SMM-based watchdog for HyperCore code



Windows®



Password

AppSpace™



ManageSpace™

HyperCore™

SecureCore BIOS

1 Virtualization-based **MALWARE**

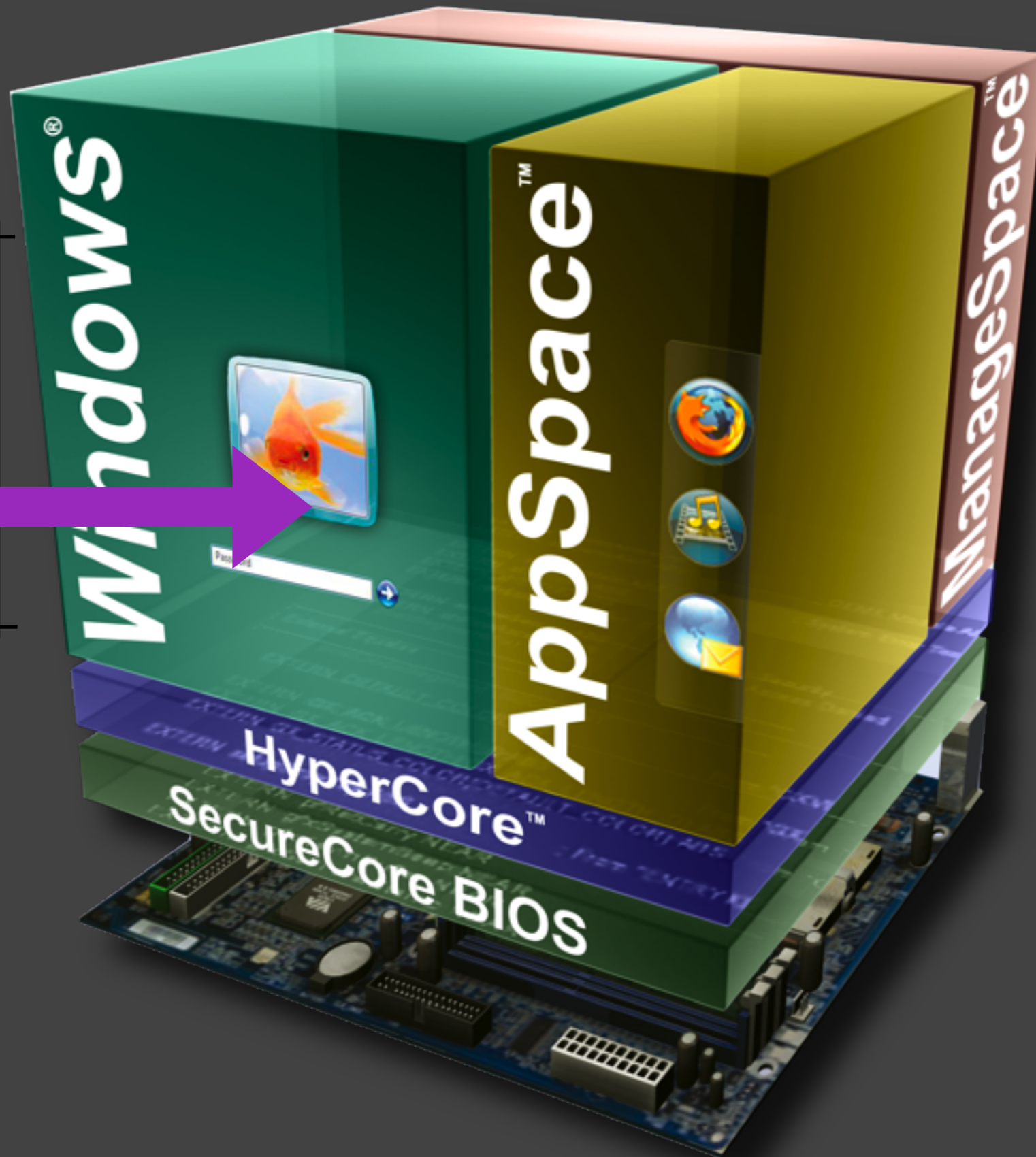
2 Using Virtual Machines for **ISOLATION**

3 **NESTED** virtualization



NESTED virtualization

What if a user wants to run e.g. Virtual PC here?



Hypervisor (Primary)

VM₁

VM₂ (Nested Hypervisor)

VM₃

VM₄

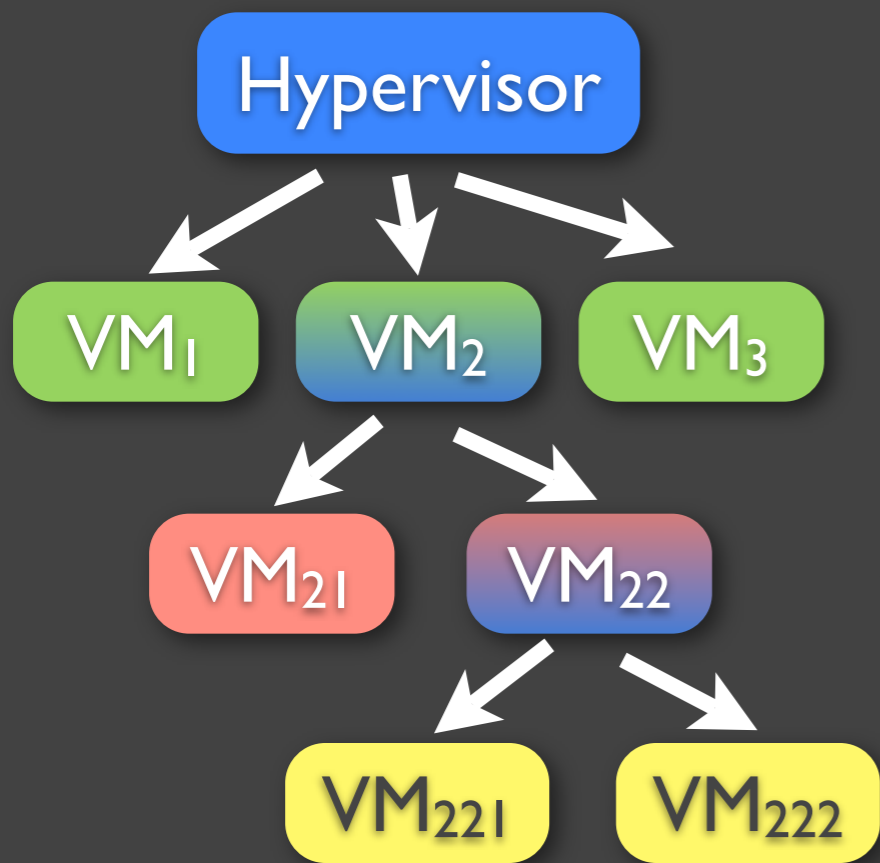
VM₂₁

VM₂₂

VM₂₂₁

VM₂₂₂

Idea of how to handle this situation...

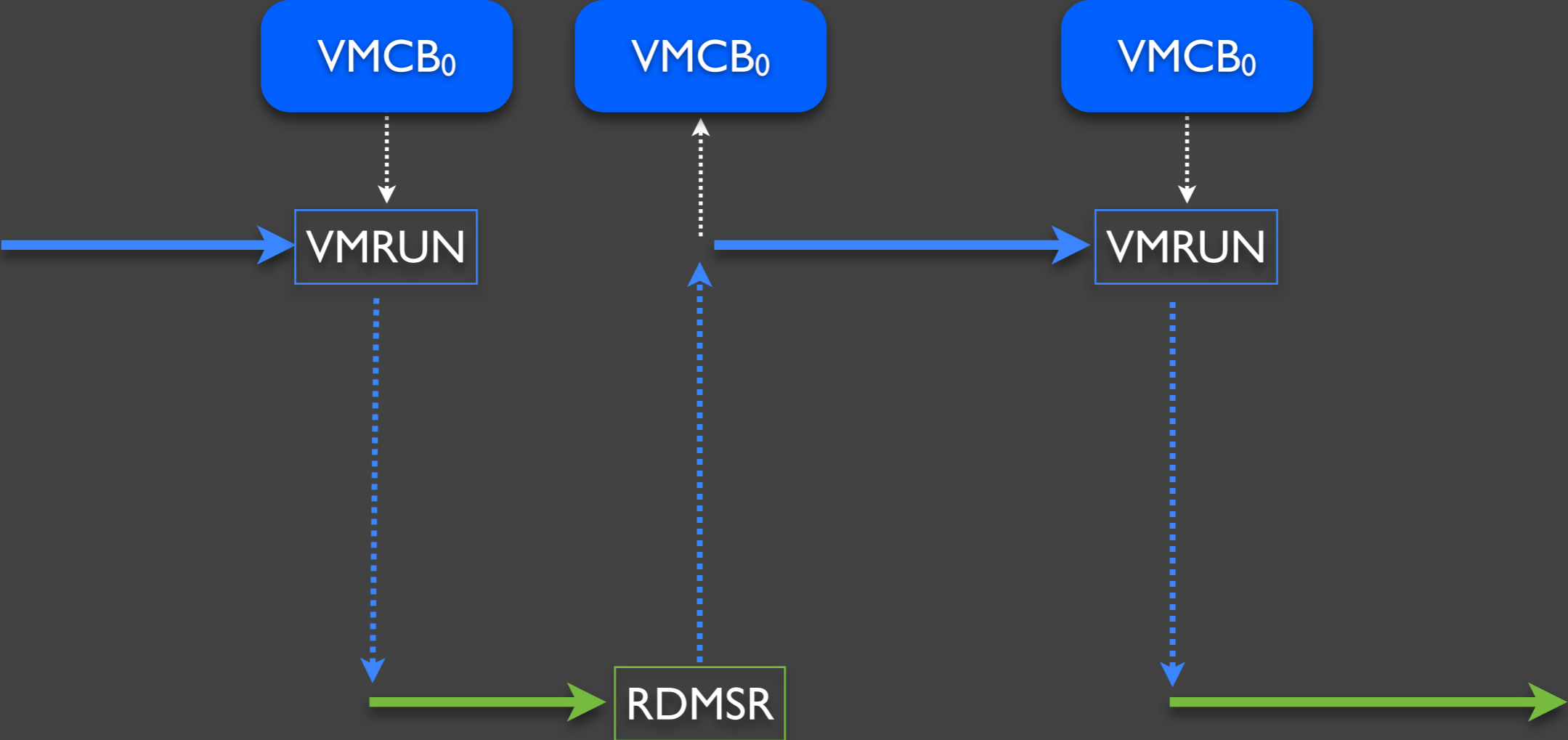


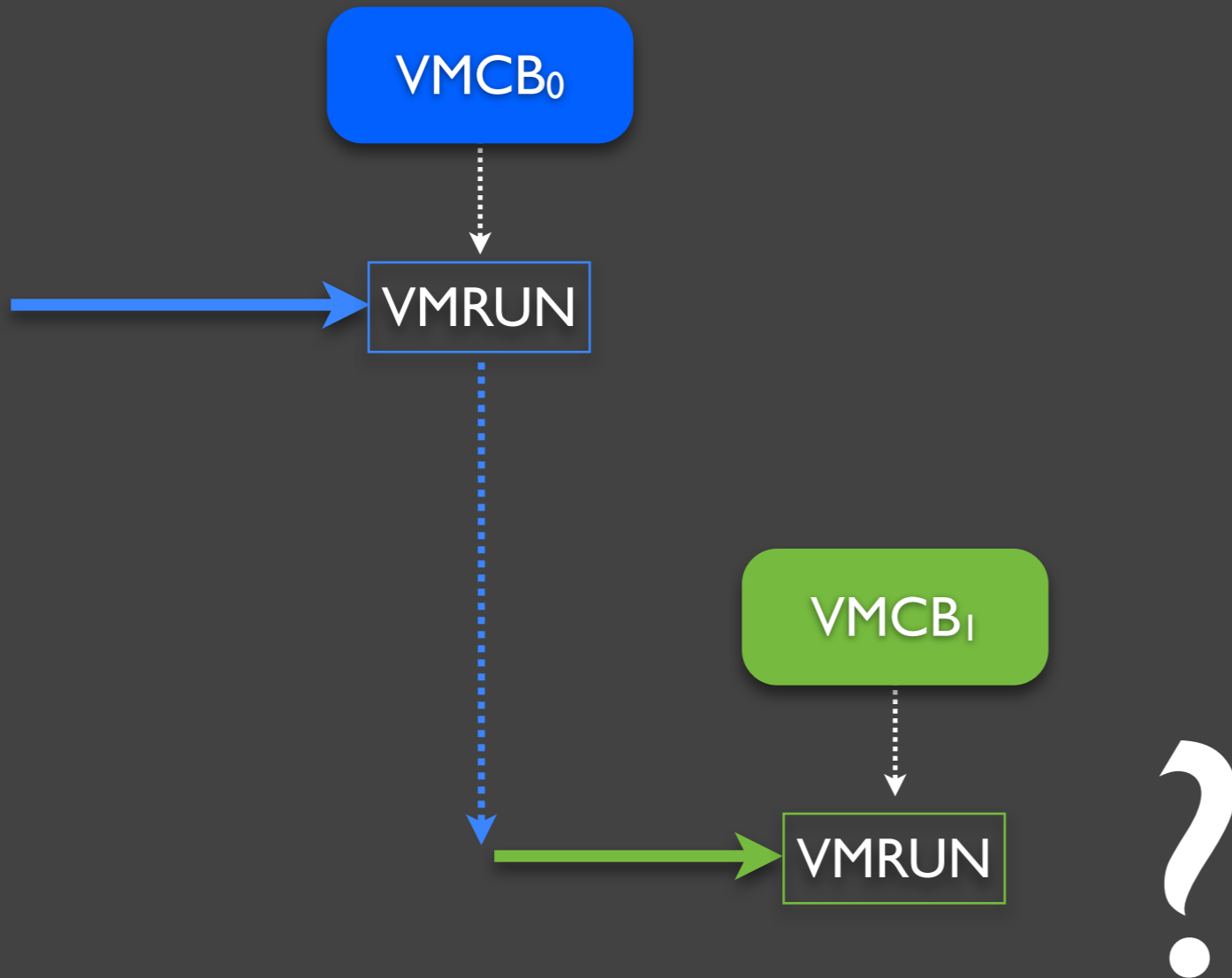
Hypervisor

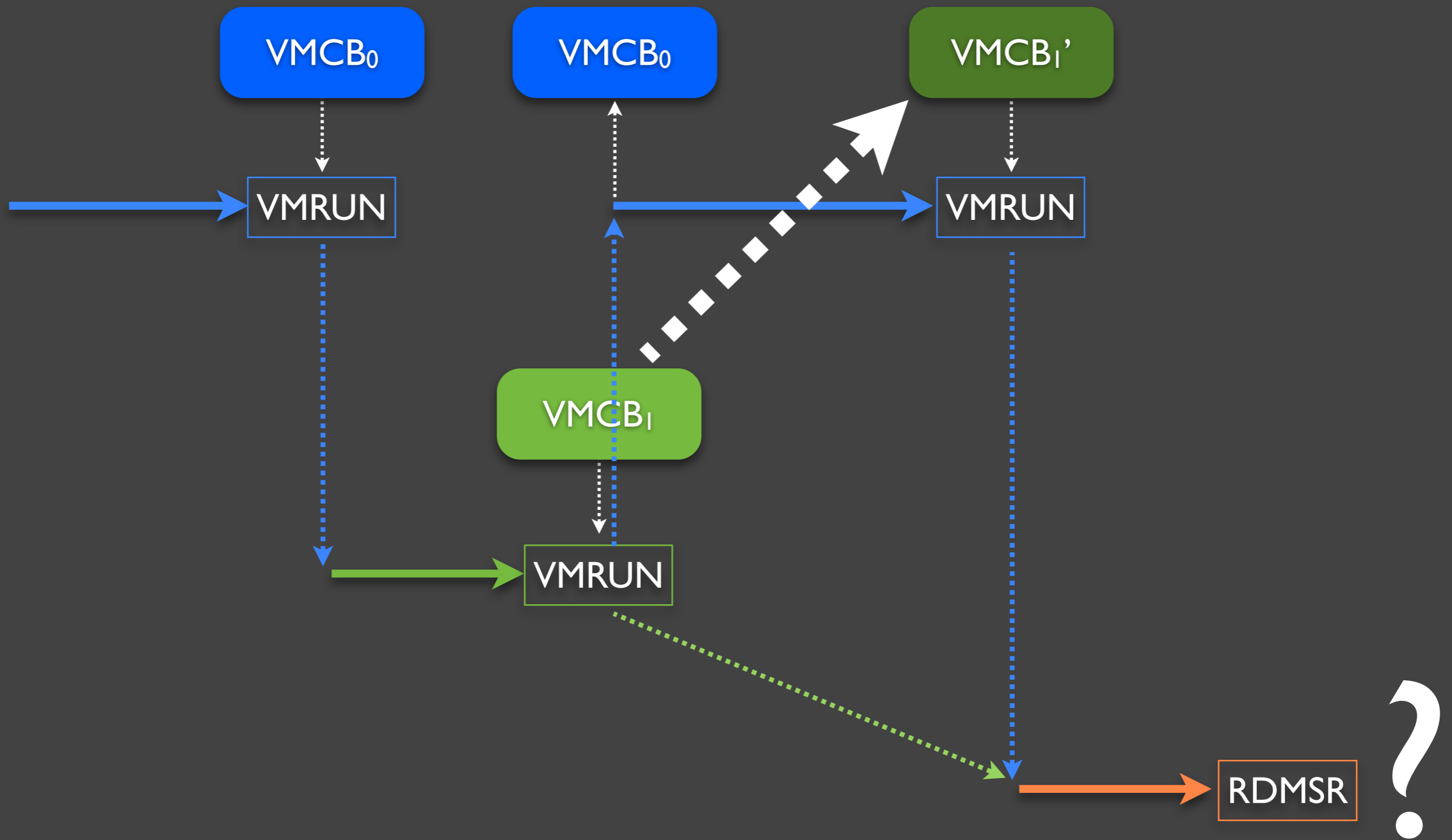


Now, lets look at the actual details :)

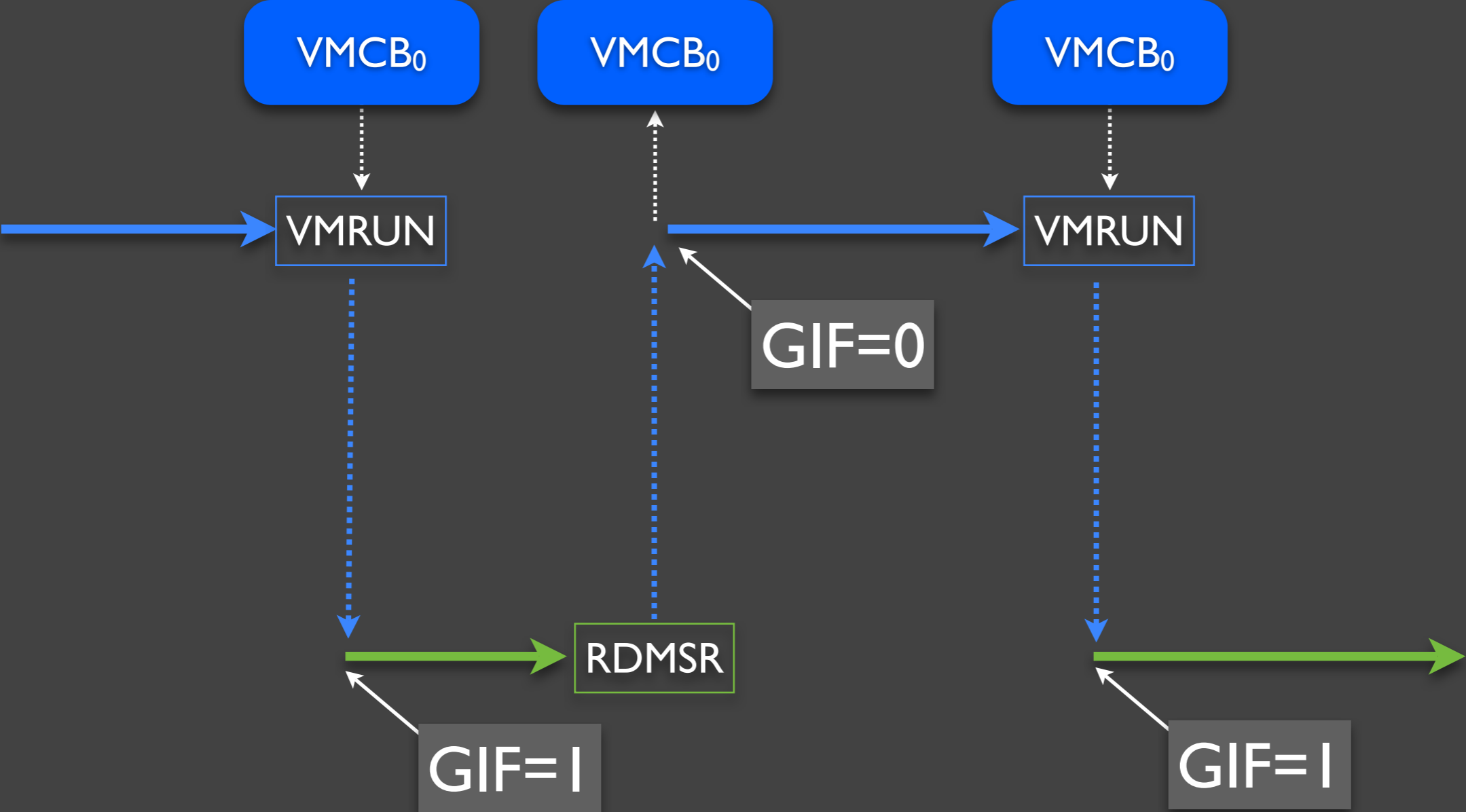
Let's start with AMD-V...

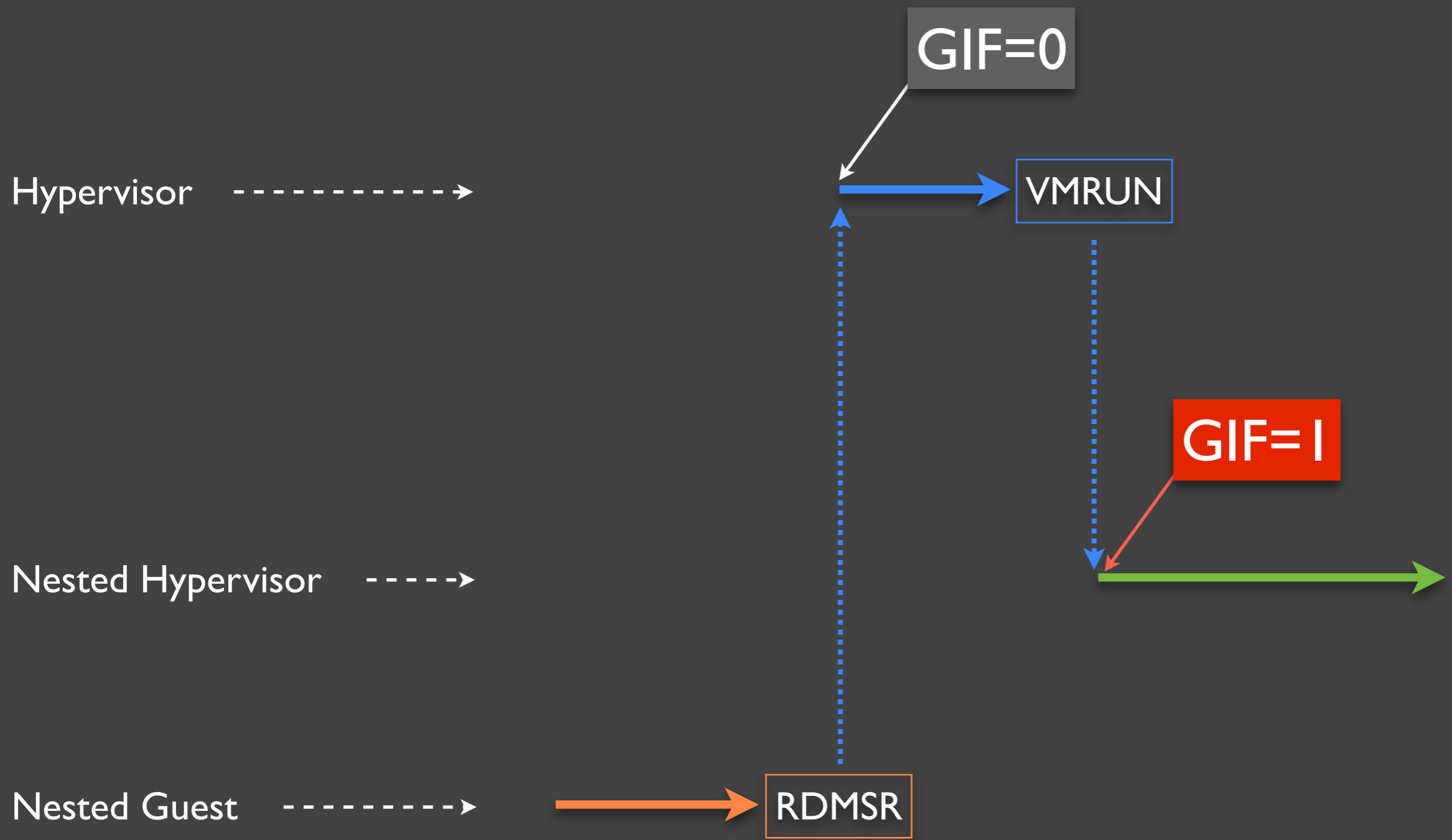






Looks convincing but won't work with more complex
hypervisors...





- Hypervisors expect to have GIF=1 when VMEXIT occurs...
- They might not be prepared to handle interrupts just after VMEXIT from guests!
- ... but when we resume the nested hypervisor CPU sets GIF=1, because we do this via VMRUN, not VMEXIT...

Getting around the “GIF Problem”

- We need to “emulate” that GIF is 0 for the nested hypervisor
- We stop this emulation when:
 - The nested hypervisor executes STGI
 - The nested hypervisor executes VMRUN
- How do we emulate it?

GIF0 emulation

- $VMCB_i.V_INTR_MASKING = 1$
- Host's $RFLAGS.IF = 0$
- Intercept NMI, SMI, INIT, #DB and held (i.e. record and reinject) or discard until we stop the emulation

Additional details

- Need to also intercept VMLOAD/VMSAVE
- Need to virtualize VM_HSAVE_PA
- ASID conflicts

Hypervisor: ASID = 0

Conflicting ASIDs!

Nested Hypervisor: ASID = 1
(but thinks that has ASID = 0)

Nested Guest: ASID = 1
(assigned by the nested hypervisor)

But we can always reassign the ASID in the VMCB “prim”
that we use to run the nested guest.

Performance Impact

- One additional #VMEXIT on every #VMEXIT that would occur in a non-nested scenario
- One additional #VMEXIT when the nested hypervisor executes: STGI, CLGI, VMLOAD, VMSAVE
- Lots of space for optimization though



```
Administrator: Command Prompt
C:\tmp\nbp-0.30>\tools\w2k_load.exe bin\amd64\newbp.sys
// w2k_load.exe
// SBS Windows 2000 Driver Loader V1.00
// 08-27-2000 Sven B. Schreiber
// sbs@orgon.com

Loading "bin\amd64\newbp.sys" ... OK

C:\tmp\nbp-0.30>\tmp\bpknock.exe 0xbabecafe
knock answer: 0x69696969

C:\tmp\nbp-0.30>"\Program Files (x86)\Microsoft Virtual PC\Virtual PC.exe"

C:\tmp\nbp-0.30>\tmp\bpknock.exe 0xbabecafe
knock answer: 0x69696969

C:\tmp\nbp-0.30>
```

Virtual PC Console

File	Action	Help
	Linux install	Not running
	Vista	Not running
	Windows XP	Running

Buttons: New..., Settings, Remove, Close...

Windows XP - Microsoft Virtual PC 2007

Action Edit CD Floppy Help

Command Prompt

```
Microsoft Windows XP [Version 5.1.2600]
(C) Copyright 1985-2001 Microsoft Corp.

C:\Documents and Settings\useruser>cd \tmp

C:\tmp>bpknock.exe 0xbabecfe
knock answer: 0

C:\tmp>bpknock.exe 0xbabecafe
knock answer: 0x69696969

C:\tmp>
```

<http://bluepillproject.org>

How AMD could help?

- AMD could add an additional field to VMCB: “EmulateGif0ForGuest”
- Additionally: virtualize STGI and CLGI when the above field is set to improve performance
- Seems simple to do: just a few additional lines in the microcode... :)

Further thinking...

- Virtualizing DEV for the nested hypervisor that makes use of DEV?
- Virtualizing IOMMU for the IOMMU-aware nested hypervisor?
- Virtualizing Nested Paging mechanism for the NP-aware nested hypervisor?

How about Intel VT-x?

Nested virtualization on VT-x

- No GIF bit - no need to emulate “GIF0” for the nested hypervisor :)
- No Tagged TLB - No ASID conflicts :)
- However:
 - VMX instructions can take memory operands - need to use complex operand parser
 - No tagged TLB - potentially bigger performance impact

Nested VT-x: Status

- We “pretty much” have that working already
- Code is messy and should be rewritten
 - e.g. the operand parser

What Intel could do?

- Extend info provided by:

`VMCS.VMX_INSTRUCTION_INFO`

So that we don't need to parse memory operand manually

- Tagged TLB for better performance
- Other optimization?

Who else does Nested (hardware-based) Virtualization?

IBM z/VM hypervisor on IBM System z™ mainframe

“Running z/VM in a virtual machine (that is, z/VM as a guest of z/VM, also known as “second-level” z/VM) is functionally supported but is intended only for testing purposes for the second-level z/VM system and its guests (called “third-level” guests).”

-- <http://www.vm.ibm.com/pubs/hcsf8b22.pdf>



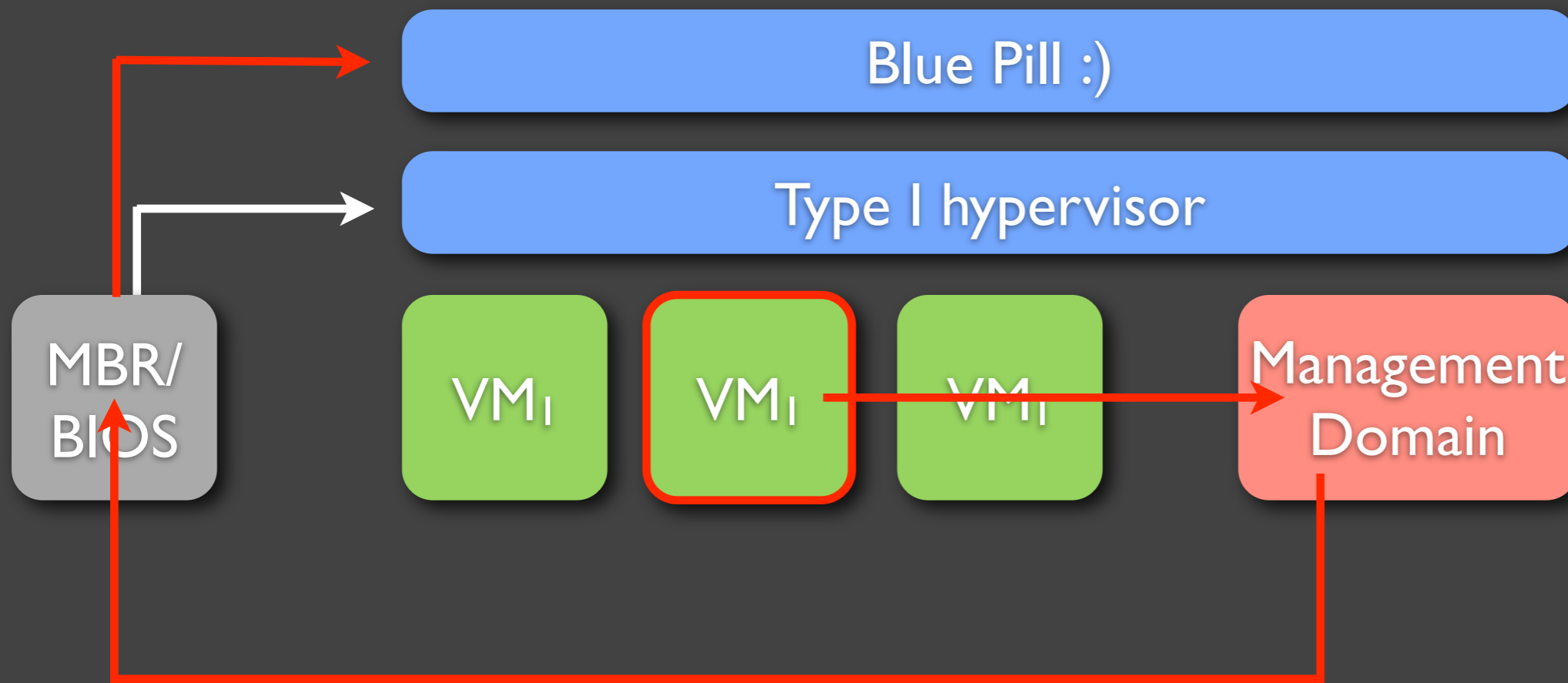
IBM System z10, source: ibm.com

Confusion

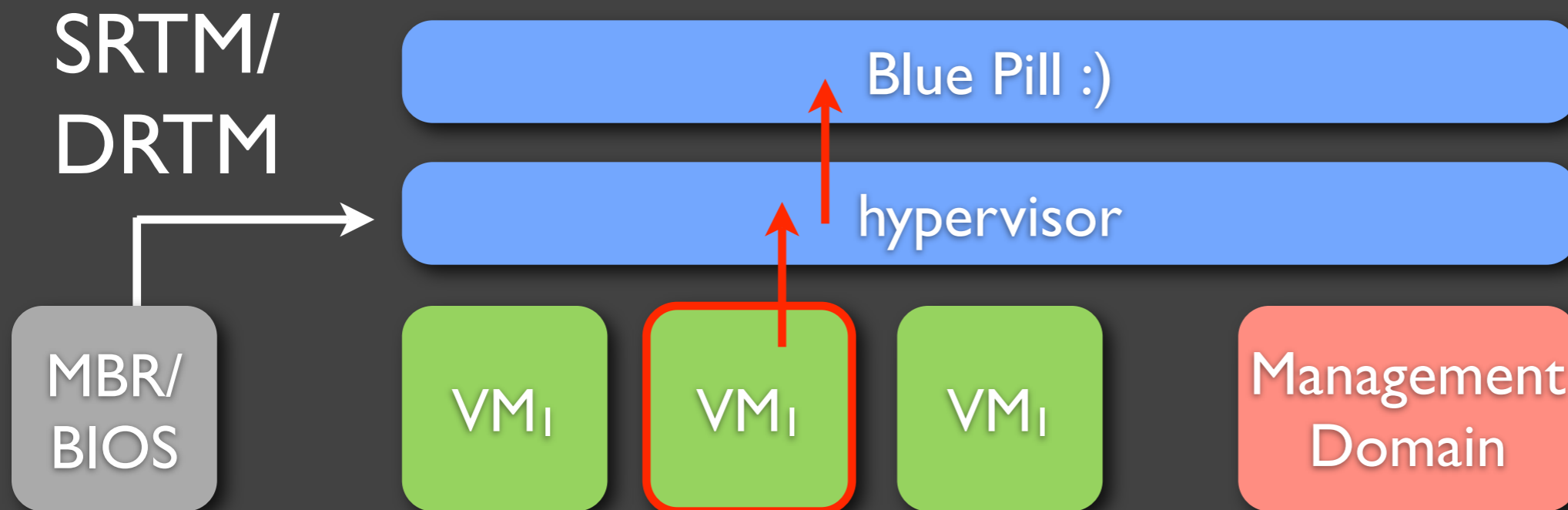
- AMD Nested Page Tables != Nested Virtualization!
- NPT is a hardware alternative to Shadow Page Tables (a good thing, BTW)
- NPT is also called: Rapid Virtualization Indexing

Nested Virtualization: Security Implications





Solution: ensure hypervisor integrity via SRTM or DRTM



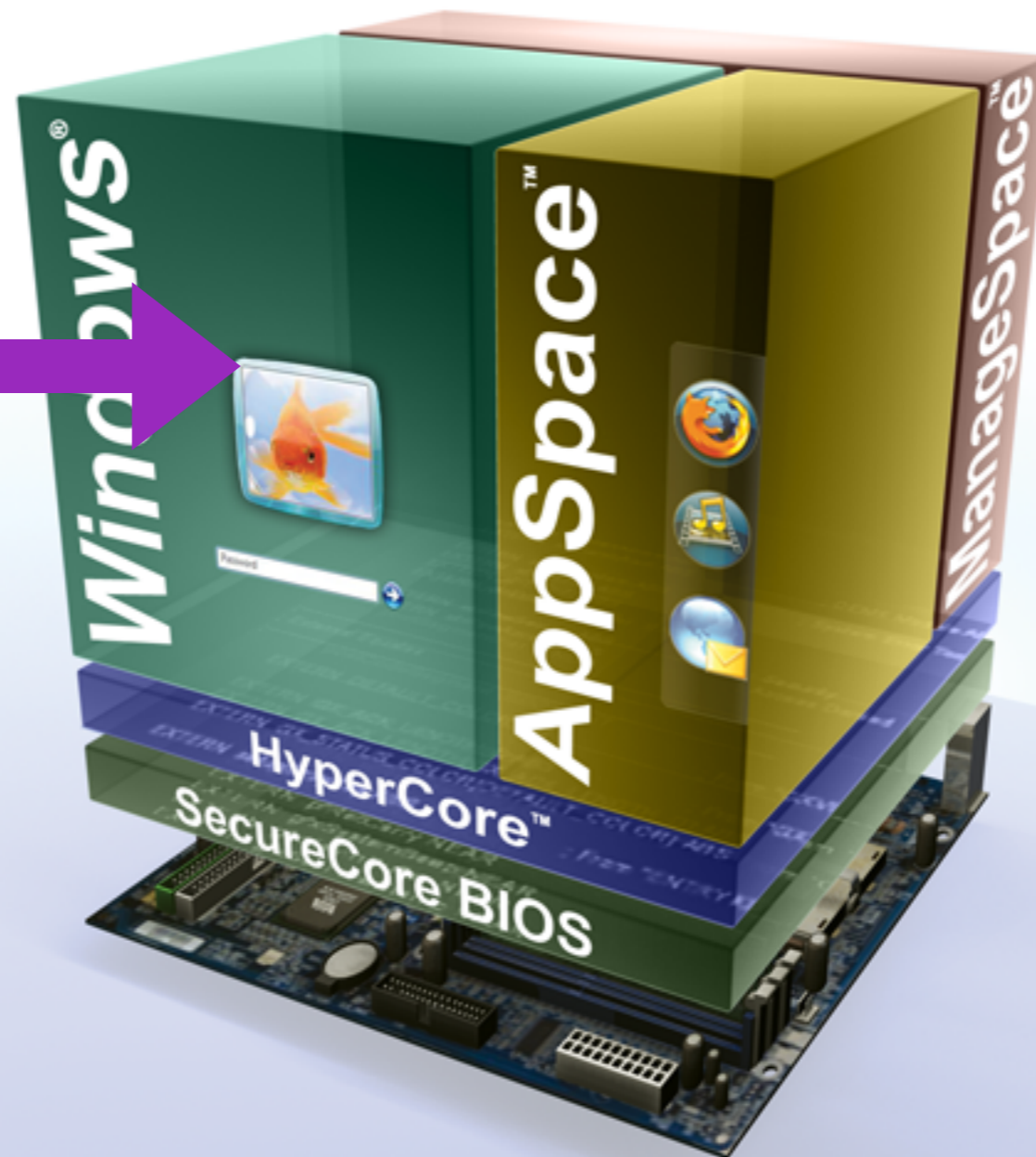
SRTM/DRTM do not protect the **already loaded hypervisor**, from being exploited **if it is buggy!**

Keep hypervisors very slim!
Do not put drivers there!

Nested Virtualization: Useful Applications

HyperSpace™

What if a user wants to run e.g. Virtual PC here?



Phoenix Technologies has supported the research on
nested hypervisors since Fall 2007

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Summary

- Virtualization technology could be used to improve security on desktop systems
- However there are non-trivial challenges in making this all working well...
- ... and not to introduce security problems instead...
- Virtualization is cool ;)

Invisible Things Lab
<http://invisiblethingslab.com>